

# Generating Combinations

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Let  $S$  be an  $n$ -set. For convenience of generating combinations, we take  $S$  to be the set

$$S = \{x_{n-1}, x_{n-2}, \dots, x_1, x_0\}.$$

Each subset  $A$  of  $S$  can be identified as a function  $\chi_A : S \rightarrow \{0, 1\}$ , called the characteristic function of  $A$ , defined by

$$\chi_A(x) = \begin{cases} 1 & \text{if } x \in A; \\ 0 & \text{if } x \notin A. \end{cases}$$

In practice,  $\chi_A$  is represented by a 0 – 1 sequence or a base 2 numeral. For example, for  $S = \{x_7, x_6, x_5, x_4, x_3, x_2, x_1, x_0\}$ ,

$\emptyset$	00000000	0
$\{x_7, x_5, x_2, x_1\}$	10100110	
$\{x_6, x_5, x_3, x_1, x_0\}$	01101011	
$\{x_7, x_6, x_5, x_4, x_3, x_2, x_1, x_0\}$	11111111	

**Algorithm** for generating combinations of  $\{x_{n-1}, x_{n-2}, \dots, x_1, x_0\}$  :

Begin with  $a_{n-1}a_{n-2} \cdots a_1a_0 = 00 \dots 00$ .

While  $a_{n-1}a_{n-2} \cdots a_1a_0 \neq 11 \dots 11$ , do

- (1) Find the smallest integer  $j$  such that  $a_j = 0$ .
- (2) Replace  $a_j$  by 1 and each of  $a_{j-1}, \dots, a_1, a_0$  by 0.

The algorithm stops when  $a_{n-1}a_{n-2} \cdots a_1a_0 = 11 \dots 11$ .

For  $n = 3$ , the algorithm produces the list

000	100
001	101
010	110
011	111.

For  $n = 4$ , the algorithm produces the list

0000	0100	1000	1100
0001	0101	1001	1101
0010	0110	1010	1110
0011	0111	1011	1111.

1. For  $n = 8$ , find the position of the combination  $\{x_6, x_5, x_3, x_1, x_0\}$  in the list.
2. Which combination comes after  $\{x_6, x_5, x_3, x_1, x_0\}$ ?

## Gray code order

The *unit  $n$ -cube*  $Q_n$  is a graph whose vertex set is the set of all 0 – 1 sequences of length  $n$ , and two sequences are adjacent if they differ in only one place.

A *Gray code of order  $n$*  is a path of  $Q_n$  that visits every vertex of  $Q_n$  exactly once, i.e., a Hamilton path of  $Q_n$ . For example,

$$000 \longrightarrow 001 \longrightarrow 101 \longrightarrow 100 \longrightarrow 110 \longrightarrow 010 \longrightarrow 011 \longrightarrow 111$$

is a Gray code of order 3. It is obvious that this Gray code can not be a part of any Hamilton cycle since 000 and 111 are not adjacent.

A cyclic Gray code of order  $n$  is a Hamilton cycle of  $Q_n$ . For example, the closed path

$$000 \longrightarrow 001 \longrightarrow 011 \longrightarrow 010 \longrightarrow 110 \longrightarrow 111 \longrightarrow 101 \longrightarrow 100 \longrightarrow 000$$

is a cyclic Gray code of order 3.

### Reflected Gray codes (inductive definition).

For  $n = 1$ , we have the Gray code  $0 \longrightarrow 1$ .

For  $n = 2$ , we use  $0 \longrightarrow 1$  to produce the path  $00 \longrightarrow 01$  by adding a 0 in the front, and use  $1 \longrightarrow 0$  to produce  $11 \longrightarrow 10$  by adding a 1 in the front, then combine the two paths to produce the Gray code  $00 \longrightarrow 01 \longrightarrow 11 \longrightarrow 10$ .

For  $n = k$ , we combine the Gray code of order  $k - 1$  with 0 added in the front of each  $(k - 1)$ -tuple, and the Gray code of order  $k - 1$  in the reverse order with 1 added in the front of each  $(k - 1)$ -tuple.

The Gray codes obtained in this way are called *reflected Gray codes*.

### Reflected Gray codes (algorithm).

**Algorithm** for generating reflected Gray codes of order  $n$ :

Begin with  $a_{n-1}a_{n-2} \cdots a_1a_0 = 00 \dots 00$ .

While  $a_{n-1}a_{n-2} \cdots a_1a_0 \neq 10 \dots 00$ , do

(1) If  $a_{n-1} + a_{n-2} + \cdots + a_1 + a_0 = 00 \dots 00 = \text{even}$ , then change  $a_0$  (from 0 to 1 or 1 to 0).

(2) If  $a_{n-1} + a_{n-2} + \cdots + a_1 + a_0 = 00 \dots 00 = \text{odd}$ , find the smallest  $j$  such that  $a_j = 1$  and change  $a_{j+1}$  (from 0 to 1 or 1 to 0).

1. Construct the reflected Gray code of order  $n = 3$  using the algorithm.
2. For  $n = 8$ , which 8-tuples follow 10100110 and 00011111 in the reflected Gray code?