## Review I

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If you put $n+1$ pigeons in $n$ pigeonholes then at least one hole would have more than one pigeon.
If $n(r-1)+1$ objects are put into $n$ boxes, then at least one of the boxes contains $r$ or more of the objects.
If the average of $n$ nonnegative integers $a_{1}, a_{2}, \ldots a_{n}$ is greater than $r-1$, i.e.,

$$
\frac{a_{1}+a_{2}+\cdots+a_{n}}{n}>r-1,
$$

then at least one of the integers is greater than or equal to $r$.
The number of $r$-permutations of an $n$-set equals

$$
P(n, r)=n(n-1) \cdots(n-r+1)=\frac{n!}{(n-r)!} .
$$

The number of permutations of an $n$-set is $P(n, n)=n!$. The number of circular $r$-permutations of an $n$-set equals

$$
\frac{P(n, r)}{r}=\frac{n!}{(n-r)!r} .
$$

The number of circular permutations of an $n$-set is equal to $(n-1)$ ! The number of $r$-combinations of an $n$-set equals

$$
\binom{n}{r}=\frac{P(n, r)}{r!}=\frac{n!}{(n-r)!r!} .
$$

The number of $r$-permutations of the multiset $\left\{\infty \cdot x_{1}, \infty \cdot x_{2}, \ldots, \infty \cdot x_{k}\right\}$ equals $k^{r}$. The number of permutations of the multiset $\left\{n_{1} \cdot x_{1}, n_{2} \cdot x_{2}, \ldots, n_{k} \cdot x_{k}\right\}$ equals

$$
\frac{n!}{n_{1}!n_{2}!\cdots n_{k}!}, \quad \text { where } n=n_{1}+n_{2}+\cdots+n_{k}
$$

Then the number of $r$-combinations of the multiset $\left\{\infty \cdot a_{1}, \infty \cdot a_{2}, \ldots, \infty \cdot a_{k}\right\}$ (the number of $r$-combinations with repetition allowed) equals $\binom{k+r-1}{r}=\binom{k+r-1}{k-1}$.
The number of nonnegative integer solutions for the equation $x_{1}+x_{2}+\cdots+x_{k}=r$ equals $\binom{k+r-1}{r}=\binom{k+r-1}{k-1}$.
The number of positive integer solutions for the equation $x_{1}+x_{2}+\cdots+x_{k}=r$ equals $\binom{r-1}{k-1}$. The number of ways to place $r$ identical balls into $k$ distinct boxes equals $\binom{k+r-1}{r}=\binom{k+r-1}{k-1}$. The number of ways to place $r$ identical balls into $k$ distinct boxes such that no box remains empty equals $\binom{r-1}{k-1}$.

Algorithm for generating the permutations of $\{1,2, \ldots, n-1, n\}$ :
Begin with $\overleftarrow{1} \overleftarrow{2} \ldots \overleftarrow{n}$.
While there exists a mobile integer, do
(1) Find the largest mobile integer $m$
(2) Switch $m$ and the adjacent integer its arrow points to.
( 30 Switch thew direction of all the arrows above integers $p$ with $p>m$.
Algorithm 1 for construction of a permutation from its inversion sequence ( $a_{1}, a_{2}, \ldots, a_{n}$ ):
(n) Write down $n$.
(n-k) Insert $n-k$ to the right of the $a_{n-k}$ th existing number

Algorithm 2 for construction of a permutation from its inversion sequence ( $a_{1}, a_{2}, \ldots, a_{n}$ ): (0) Mark down $n$ empty spaces.

For $k=1$ till $n$
Put $k$ into the $a_{k}+1$ st empty space from the left.
Algorithm for generating combinations of $\left\{x_{n-1}, x_{n-2}, \ldots, x_{1}, x_{0}\right\}$ :
Begin with $a_{n-1} a_{n-2} \cdots a_{1} a_{0}=00 \ldots 00$.
While $a_{n-1} a_{n-2} \cdots a_{1} a_{0} \neq 11 \ldots 11$, do
(1) Find the smallest integer $j$ such that $a_{j}=0$.
(2) Replace $a_{j}$ by 1 and each of $a_{j-1}, \ldots, a_{1}, a_{0}$ by 0 .

The algorithm stops when $a_{n-1} a_{n-2} \cdots a_{1} a_{0}=11 \ldots 11$.
Algorithm for generating reflected Gray codes of order $n$ :
Begin with $a_{n-1} a_{n-2} \cdots a_{1} a_{0}=00 \ldots 00$.
While $a_{n-1} a_{n-2} \cdots a_{1} a_{0} \neq 10 \ldots 00$, do
(1) If $a_{n-1}+a_{n-2}+\cdots+a_{1}+a_{0}=$ even, then change $a_{0}$ (from 0 to 1 or 1 to 0 ).
(2) If $a_{n-1}+a_{n-2}+\cdots+a_{1}+a_{0}=$ odd, find the smallest $j$ such that $a_{j}=1$ and change $a_{j+1}$ (from 0 to 1 or 1 to 0 ).
Algorithm for generating $r$-combinations of $S=\{1,2, \ldots, n-1, n\}$ :
Begin with $12 \cdots r$.
While $a_{1} a_{2} \cdots a_{r} \neq(n-r+1) \cdots(n-1) n$, do
(1) Find the largest integer $k$ such that $a_{k}<n$ and $a_{k}+1$ is not in the $a_{1} a_{2} \cdots a_{r}$.
(2) Replace $a_{1} a_{2} \cdots a_{r}$ with

$$
a_{1} a_{2} \cdots a_{k-1}\left(a_{k}+1\right)\left(a_{k}+2\right) \cdots\left(a_{k}+r-k+1\right)
$$

Algorithm for a linear extension of an $n$-poset:
Step 1. Choose a minimal element $x_{1}$ from X (with respect to the ordering $\leq$ ).
Step 2. Delete $x_{1}$ from $X$; choose a minimal element $x_{2}$ from $X$.
Step 3. Delete $x_{2}$ from $X$ and choose a minimal element $x_{3}$ from $X$.
...
Step n. Delete $x_{n-1}$ from $X$ and choose the only element $x_{n}$ in $X$.

